

# 6.5630 Advanced Topics in Cryptography

## Problem Set 3

### Optional PSET, No Due Date

This problem set has a single problem.

**Breaking Functional Encryption.** Recall the construction of a predicate encryption scheme supporting the *orthogonality predicate* (alternatively, linear functions with equality test) from the LWE assumption (see Lecture 8 notes, section 3). Such a function is defined by a vector  $y \in \mathbb{Z}_q^L$ , takes as input  $x \in \mathbb{Z}_q^L$  and outputs 1 if

$$\langle y, x \rangle = 0 \pmod{q}$$

We showed in class that this construction is secure as a weakly hiding predicate encryption. That is, an adversary who gets secret keys corresponding to a number of vectors (linear functions)  $y_1, \dots, y_k$  for a polynomially large  $k$  such that

$$\langle x, y_i \rangle \neq 0 \pmod{q} \text{ for all } i,$$

cannot learn anything about  $x$ . Your task is to break this scheme when playing the strong hiding game. That is, when the challenge is  $(x, x')$  such that  $\langle x, y_i \rangle = \langle x', y_i \rangle \pmod{q}$  for all  $i$  (but some of the inner products could be 0).

Concretely, your task is to design queries  $y_i$  such that the outputs  $\langle x, y_i \rangle \pmod{q}$  does not reveal  $x$ , yet a ciphertext  $\text{Enc}(\text{mpk}, x)$  together with the secret keys  $\text{sk}_{y_1}, \dots, \text{sk}_{y_\ell}$  completely reveals  $x$ .